

CYLINDRIC ALGEBRAS OF FIRST-ORDER LANGUAGES

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ABSTRACT. We show when two countable first-order languages have isomorphic cylindric algebras.

Introduction. The *cylindric algebra of a language* is the ω -dimensional cylindric algebra of formulas of the language modulo logical equivalence. We classify first-order languages according to the isomorphism types of their cylindric algebras. For two languages involving only predicate symbols, their cylindric algebras are isomorphic iff for every $n \in \omega$, both languages have the same number of predicate symbols with at least n places. This solves the classification problem for free ω -dimensional locally-finite cylindric algebras [4, Problem 2.8, p. 463] since these algebras are exactly the cylindric algebras of languages involving only predicate symbols.

For any first-order theory, the topological space of its models with EC_Δ classes (classes which consist of all models of some theory) as closed sets is a natural dual for the Lindenbaum-Tarski algebra of the theory. By enriching the class of models to the concrete category of models and isomorphisms and replacing the topology with the ultraproduct operations, we obtain a natural model-theoretic dual for the cylindric algebra of formulas of the theory. This duality is used in demonstrating that various properties of languages are invariant under cylindric isomorphism.

The basic tool used in isomorphism construction is a decomposition version of the Cantor back-and-forth construction [7].

Notation. All languages, theories and sets of symbols will be countable. All theories will be first-order theories with equality. "Symbol" will mean nonlogical symbol. The *type* of an n -ary predicate symbol is n ; that of an n -ary operation symbol is $n + \frac{1}{2}$. Propositional letters are regarded as 0-ary predicate symbols and constants as 0-ary operation symbols. Let $\frac{1}{2}\omega = \{\frac{1}{2}, 1\frac{1}{2}, 2\frac{1}{2}, \dots\}$. The *(full) type sequence* of a language is the sequence $\langle a_0, a_{\frac{1}{2}}, a_1, a_{1\frac{1}{2}}, \dots \rangle$ where a_i is the number (possibly ω) of symbols of the language of type i . The *reduced type sequence* is the sequence $\langle b_0, b_{\frac{1}{2}}, b_1, b_{1\frac{1}{2}}, \dots \rangle$ where b_i is ω if there are

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infinitely many symbols of type $\geq i$ and, if not, the number of symbols of type i .

The *unboundedness* of the language is the largest i , if any, such that $b_i = \omega$; it is undefined if no b_i is ω and ω if every b_i is ω .

We shall occasionally wish to consider symbols other than equality as quasi-logical built-in symbols. Let S be a fixed set of such symbols, let Σ be a set of axioms on S and let $s = \langle S, \Sigma \rangle$. For any set R of symbols, let $L_s(R)$ be the language whose set of symbols is $S \cup R$ and let $Th_s(R)$ be the theory whose language is $L_s(R)$ and whose axioms are Σ . We write $L(R)$ and $Th(R)$ for $L_s(R)$ and $Th_s(R)$ when $s = \langle \emptyset, \emptyset \rangle$. Suppose T is a theory whose set of symbols includes S . If R' is a subset of T 's set of symbols, let $T|R'$ be the restriction of T to the sentences of $L_s(R')$. If σ is a sentence of T 's language, let $T\sigma$ be the theory obtained by adding σ to the axioms of T . Let $cyl_s(T)$ be the cylindric algebra of formula's of T 's language modulo equivalence with respect to T plus constants a_S for $S \in S$ where a_S is $Sx_1 \dots x_n$ if S is an n -ary predicate symbol and $Sx_1 \dots x_n = x_{n+1}$ if S is an n -ary operation symbol. Let $\cong_s$ be the relation on theories such that $T \cong_s T'$ iff $cyl_s(T) \cong cyl_s(T')$. We write $\cong$ for $\cong_s$ when $s = \langle \emptyset, \emptyset \rangle$. If $T \cong T'$, T and T' are *cylindrically isomorphic*. Being cylindrically isomorphic is equivalent to being bilaterally interpretable [6], being synonymous [1] or, if there are no constant or operation symbols, being related [8].

Cylindric duality. In this section and the next, the S and Σ of the previous section will be empty. Let Cyl be the category consisting of the algebras $cyl(T)$, T a theory, and cylindric homomorphisms. Given a theory T , an ultraproduct operation on the class of its models is one which, for some index set I and ultrafilter D , sends an I -indexed family $\langle \mathcal{U}_i : i \in I \rangle$ of models of T to the ultraproduct $\Pi_D \mathcal{U}_i$. Let $mod\text{-}ult(T)$ be the concrete category of models of T and isomorphisms operated on by the ultraproduct operations. Let $Mod\text{-}Ult$ be the category whose objects are the categories $mod\text{-}ult(T)$, T a theory, and whose morphisms are functors $h: mod\text{-}ult(T) \rightarrow mod\text{-}ult(T')$ such that $h \circ$ (the forgetful functor) = the identity and h commutes with all ultraproduct operations.

THEOREM 1 (DUALITY THEOREM). *The categories Cyl and $Mod\text{-}Ult$ are dual.*

PROOF. For any theory T , let $cyl(T)^* = mod\text{-}ult(T)$. Let T and T' be theories and suppose for convenience that their symbols are the binary predicates R and S respectively. For any $f: cyl(T') \rightarrow cyl(T)$, let $f^*: mod\text{-}ult(T) \rightarrow mod\text{-}ult(T')$ be the unique concrete functor such that for any model $\mathcal{U} = \langle X, R^{\mathcal{U}} \rangle$ of T , $f^*(\langle X, R^{\mathcal{U}} \rangle) = \langle X, \varphi^{\mathcal{U}} \rangle$ where $\varphi(x, y) = f(Sxy)$ and $\varphi^{\mathcal{U}}$ is the interpretation of φ in $\mathcal{U}$. The structure $\langle X, \varphi^{\mathcal{U}} \rangle$ is a model of T' since if $T' \models \sigma$ then, since f is a cylindric homomorphism, $T \models f(\sigma)$ and hence $\langle X, R^{\mathcal{U}} \rangle \models f(\sigma)$ and hence, since $f(\sigma)$ is the result of replacing S with φ in σ , $\langle X, \varphi^{\mathcal{U}} \rangle \models \sigma$. Clearly, f^* preserves ultraproducts.

Easily, if i is an identity map of Cyl , then i^* is an identity map of Mod-Ult and if f and g are maps of Cyl , then $(f \circ g)^* = g^* \circ f^*$. Hence $(\)^*$ is a contravariant functor from Cyl to Mod-Ult .

For any theory T let $\text{mod-ult}(T)^* = \text{cyl}(T)$. Again suppose T and T' are theories whose symbols are the binary predicates R and S respectively. Let $h: \text{mod-ult}(T) \rightarrow \text{mod-ult}(T')$ be a functor in Mod-Ult . Then the map U_h defined by $U_h(\mathfrak{U}) = \langle X, R^{\mathfrak{U}}, S^h(\mathfrak{U}) \rangle$ for $\mathfrak{U} = \langle X, R^{\mathfrak{U}} \rangle$ commutes with the ultraproduct operations and preserves isomorphisms. Hence by Kochen [6, Theorem 12.1] and Shelah [9] (the latter shows that the ultralimit condition of the former may be omitted) there is a formula $\varphi(x, y)$ involving only R such that for any $\mathfrak{U} = \langle X, R^{\mathfrak{U}} \rangle$, $\varphi^{\mathfrak{U}} = S^h(\mathfrak{U})$. Clearly this $\varphi(x, y)$ is unique modulo equivalence with respect to T . Let $f: \text{cyl}(L(\{S\})) \rightarrow \text{cyl}(T)$ be the unique homomorphism such that $f(Sxy) = \varphi(x, y)$ and let $p: \text{cyl}(L(\{S\})) \rightarrow \text{cyl}(T')$ be the homomorphism such that $p(Sxy) = Sxy$. For any sentence σ of $L(\{S\})$, if $T' \models \sigma$ and $\mathfrak{U} \models T$ then $h(\mathfrak{U}) \models T'$ and so $h(\mathfrak{U}) \models \sigma$ and so $\mathfrak{U} \models f(\sigma)$. Hence $T' \models \sigma$ implies $T \models f(\sigma)$ and so f factors through p and, since p is onto, the factorization is unique. Let h^* be the unique homomorphism such that $f = h^* \circ p$. Easily, for any identity map i of Mod-Ult , i^* is an identity map of Cyl and for any h and g , $(h \circ g)^* = g^* \circ h^*$. Finally it is clear that $\text{cyl}(T)^{**} = \text{cyl}(T)$ and $\text{mod-ult}(T)^{**} = \text{mod-ult}(T)$. It is also straightforward to verify that for any map h in Mod-Ult , $h^{**} = h$. For any map f in Cyl , $f^{**} = f$ follows from the uniqueness modulo equivalence with respect to T of the $\varphi(x, y)$ in the definition above.

COROLLARY 2. *For any theories T and T' , if $\text{cyl}(T) \cong \text{cyl}(T')$, then*

- (1) *the number of models of T with a given universe equals the number of models of T' with that universe,*
- (2) *the number of nonisomorphic models of T of a given cardinality equals the number of nonisomorphic models of T' of that cardinality, and*
- (3) *every automorphism group of a model of T is also (identical with not just isomorphic to) an automorphism group of a model of T' and vice versa.*

Our model-theoretic dual should be compared to the sheaf-theoretic duals of S. D. Comer [2] which are part model-theoretic (the base space is the topological space of models) and part algebraic (the stalks are cylindric algebras).

Invariants. An invariant of a language is a property of or a function on languages such that languages with isomorphic cylindric algebras are the same with respect to the property or are assigned the same value by the function.

Although the full type sequence is not an invariant, we have

LEMMA 3. *The reduced type sequence is an invariant.*

PROOF. The proof is basically the counting argument of Pigozzi in [4,

Theorem 2.5.51]. We show that languages with different reduced type sequences are not cylindrically isomorphic. Let T and T' be the theories of languages L and L' whose sets of symbols are R and R' and whose reduced type sequences are $\alpha = \langle a_0, a_{1/2}, a_1, \dots \rangle$ and $\beta = \langle b_0, b_{1/2}, b_1, \dots \rangle$. If $\alpha \neq \beta$, then, since they are reduced type sequences, one of α and β must be smaller reverse lexicographically. Suppose for some i , $a_i > b_i$ and $a_j \geq b_j$ for $j > i$. Thus $b_i < \omega$ and so $b_j < \omega$ for $j > i$ and for some $k \geq i$, $b_j = 0$ for $j \geq k$. Suppose for reductio ad absurdum that $f: \text{cyl}(T) \cong \text{cyl}(T')$. Let R_0 be a subset of R consisting of $b_i + 1$ symbols of type i and b_j symbols of type j for $j \in (i, k]$. Then f maps $\text{cyl}(T|R_0)$ 1-1 into $\text{cyl}(T'|R'_0)$ for some finite $R'_0 \subseteq R'$ consisting of b'_i symbols of type i for $i \in [0, k]$ where $b'_i \leq b_i$. Hence f^* , the dual of f defined in the Duality section, is a universe preserving map from the structures of $L(R'_0)$ to those of $L(R_0)$. If $\mathcal{U}$ is a structure of $L(R_0)$ with finite universe and σ a sentence of $\text{Cyl}(T|R_0)$ whose model class is $\mathcal{U}$'s isomorphism type, then $\mathcal{U}$ not in the image of f^* implies $f(\sigma) = 0$ which is impossible since f is 1-1. Hence f^* maps the finite structures of $L(R'_0)$ onto the finite structures of $L(R_0)$. Let x be a finite cardinal and let $N_0(x)$ and $N'_0(x)$ be the number of $L(R_0)$ -structures and $L(R'_0)$ -structures, respectively, over a given universe of cardinality x . Thus $N_0(x) \leq N'_0(x)$. For illustrative purposes, suppose $i \in \omega$ and $k \in \frac{1}{2}\omega$. Then

$$N_0(x) = (2^{(x^i)})^{b_i+1} (x^{(x^i)})^{b_{i+1/2}} (2^{(x^{i+1})})^{b_{i+1}} \dots (x^{(x^{k-1/2})})^{b_k}, \quad \text{and}$$

$$N'_0(x) = (2^{(x^0)})^{b'_0} (x^{(x^0)})^{b'_{1/2}} \dots (2^{(x^i)})^{b'_i} \dots (x^{(x^{k-1/2})})^{b'_k}.$$

Taking logarithms we have

$$\log_2 N_0(x) = (b_i + 1)x^i + b_{i+1/2}x^i \log_2 x + b_{i+1}x^{i+1} + \dots + b_k x^{k-1/2} \log_2 x,$$

and

$$\log_2 N'_0(x) = b'_0 x^0 + b'_{1/2} x^0 \log_2 x + \dots + b'_i x^i + \dots + b'_k x^{k-1/2} \log_2 x.$$

Since $b_i + 1 > b'_i$ and $b_j \geq b'_j$ for $j \in (i, k]$, a simple rate of growth argument shows that $\log_2 N_0(x) > \log_2 N'_0(x)$ for sufficiently large x . But this contradicts $N_0(x) \leq N'_0(x)$. Hence T and T' are not cylindrically isomorphic.

LEMMA 4. *The existence of a constant is an invariant.*

PROOF. Any structure of a language with a constant has a point which is left fixed by all automorphisms. Any language without a constant has a structure with no such fixed point. Hence no language with a constant is cylindrically isomorphic to one without a constant.

LEMMA 5. *The number of predicate symbols is an invariant.*

PROOF. The number of isomorphism types of one element structures of a language with m predicate symbols is 2^m and, by Corollary 2, this number is an invariant.

LEMMA 6. *The number of unary predicate symbols is an invariant for the class of languages whose symbols are of type 0, 1 or $1\frac{1}{2}$.*

PROOF. Suppose L and L' are languages all of whose symbols are of type 0, 1 or $1\frac{1}{2}$; suppose L has k unary predicate symbols, $k \in \omega \cup \{\omega\}$; and suppose L' has k' unary predicate symbols, $k' < k$. By interpreting the function symbols to be the identity and by choosing appropriate interpretations for the unary predicates, we can construct an L -structure with an infinite universe X which can be partitioned into 2^k infinite orbits such that the restriction of the automorphism group of the structure to an orbit is the full symmetric group of that orbit. In any L' -structure over X with the same automorphism group all the function symbols must be interpreted as identities since if $f(x) \neq x$ we can find an automorphism which moves $f(x)$ but not x , an impossibility. But any such L' -structure can be partitioned into at most $2^{k'}$ orbits. Hence $k' < k$ is impossible.

Isomorphism construction. Given an index set S , let $\mathbf{C}$ be the category of countable ω -dimensional locally-finite cylindric algebras with an S -indexed list of constants and homomorphisms which preserve the constants. The *minimal* subalgebra of an algebra in $\mathbf{C}$ is the subalgebra generated by the constants. Let $\oplus: \mathbf{C} \times \mathbf{C} \rightarrow \mathbf{C}$ be the cartesian product functor and let $\otimes: \mathbf{C} \times \mathbf{C} \rightarrow \mathbf{C}$ be the tensor product functor. For objects $\mathcal{U}$ and $\mathcal{B}$ of $\mathbf{C}$, $\mathcal{U} \otimes \mathcal{B}$ is the categorical direct sum of $\mathcal{U}$ and $\mathcal{B}$ in $\mathbf{C}$. If there are no additional constants, $\mathcal{U} \otimes \mathcal{B}$ is the free product of the cylindric algebras $\mathcal{U}$ and $\mathcal{B}$. If there are additional constants and $\mathcal{U}$ and $\mathcal{B}$ have the same minimal algebra $\mathcal{C}$, $\mathcal{U} \otimes \mathcal{B}$ is the free product of $\mathcal{U}$ and $\mathcal{B}$ amalgamated over $\mathcal{C}$ with constants interpreted as in $\mathcal{C}$. Let Φ be the class of all polynomials in $\otimes$ and $\oplus$ with coefficients in $\mathbf{C}$. This is the smallest class of functors closed under composition and containing $\otimes$, $\oplus$, the identity and projection functors, and all zero-ary functors whose unique value is in $\mathbf{C}$.

If $\mathcal{D}$ is the minimal subalgebra of $\mathcal{U}$, $\mathcal{D} \rightarrow \mathcal{U}$ will always be the canonical injection. If maps $\mathcal{U}_1 \rightarrow \mathcal{B}_1, \dots, \mathcal{U}_n \rightarrow \mathcal{B}_n$ have been given, $\varphi(\mathcal{U}_1, \dots, \mathcal{U}_n) \rightarrow \varphi(\mathcal{B}_1, \dots, \mathcal{B}_n)$ will always be the map $\varphi(\mathcal{U}_1 \rightarrow \mathcal{B}_1, \dots, \mathcal{U}_n \rightarrow \mathcal{B}_n)$.

A symmetric binary relation Q on the algebras of $\mathbf{C}$ is Φ -decomposable iff for any algebras $\mathcal{U}$ and $\mathcal{B}$ of $\mathbf{C}$, $(D1) \mathcal{U} Q \mathcal{B}$ implies $\mathcal{U}$ and $\mathcal{B}$ have isomorphic minimal subalgebras and $(D2) \mathcal{U} Q \mathcal{B}$ and $\mathcal{C} \rightarrow \mathcal{U}$ a monomorphism from a finitely generated algebra of $\mathbf{C}$ implies there is a functor $\varphi(X_1, \dots, X_n)$ in Φ , sequences $\bar{\mathcal{U}} = \langle \mathcal{U}_1, \dots, \mathcal{U}_n \rangle$ and $\bar{\mathcal{B}} = \langle \mathcal{B}_1, \dots, \mathcal{B}_n \rangle$ of algebras in $\mathbf{C}$ and isomorphisms $\varphi(\mathcal{U}_1, \dots, \mathcal{U}_n) \cong \mathcal{U}$ and $\varphi(\mathcal{B}_1, \dots, \mathcal{B}_n) \cong \mathcal{B}$ such that $\mathcal{U}_1 Q \mathcal{B}_1, \dots, \mathcal{U}_n Q \mathcal{B}_n$, and $\mathcal{C} \rightarrow \mathcal{U}$ factors through $\varphi(\mathcal{D}_1, \dots, \mathcal{D}_n) \rightarrow \varphi(\mathcal{U}_1, \dots, \mathcal{U}_n) \cong \mathcal{U}$ where $\mathcal{D}_i$ is the minimal subalgebra of $\mathcal{U}_i$.

LEMMA 7. If Q is a Φ -decomposable relation on C and $\mathcal{U}$ and $\mathcal{B}$ algebras of C , then $\mathcal{U}Q\mathcal{B}$ implies $\mathcal{U} \cong \mathcal{B}$.

PROOF. This lemma is a special case of Theorems 2 and 3 of [7].

Let S be the set of quasi-logical symbols. For any n -place predicate symbol R , let " R is trivial" be the sentence $(\forall x_1 \cdots x_n)(\neg Rx_1 \cdots x_n)$ and let " R is nontrivial" be its negation. For any n -place operation symbol R , $n \geq 1$, let " R is trivial" be $(\forall x_1 \cdots x_n)(Rx_1 \cdots x_n = x_1)$. If R is a constant symbol and a is a given constant of S , let " R is trivial" be $(R = a)$. In any case, for any theory T whose symbols include R and those of S , R is definable from S in T (" R is trivial").

The peculiar form of the next theorem is motivated by subsequent applications. Note that for any theory whose set of symbols includes S , $T|\emptyset = T|S$.

THEOREM 8. (CONSTRUCTION THEOREM). Let S be a fixed set of symbols let Σ be a set of axioms of S , and let $s = \langle S, \Sigma \rangle$. Suppose that Q is a symmetric relation on theories and that for any theories T and T' satisfying TQT' , there are sets of symbols R and R' which include S and for which $T = Th_s(R)$ and $T' = Th_s(R')$ and the following hold:

(C1). $T|\emptyset = T'|\emptyset$.

(C2). $R_1, \dots, R_n \in R$ implies there are $R_1, \dots, R_n \subseteq R$ and $R'_1, \dots, R'_n \subseteq R'$ such that $R_i \in R_i$, $R_i \cap R_j = R'_i \cap R'_j = S$ for $i \neq j$, $R = R_1 \cup \dots \cup R_n$, $R' = R'_1 \cup \dots \cup R'_n$, and $T|R_iQT'|R'_i$.

(C3). For a finite subset of R and F' a finite subset of R' implies $T|(R \sim F)QT|(R' \sim F')$.

(C4). R and R' have no constants if S has none.

(C5). If R is an n -ary predicate (operation) symbol of R such that no symbol of the same type occurs in R' , then there are symbols $R_1, \dots, R_m$ of R of types which do occur in R' and a sentence denoted by " $R_1, \dots, R_m$ codes an n -ary relation (operation)" in $L_s(\{R_1, \dots, R_m\})$ such that $T|\{R\} \cong_s T|\{R_1, \dots, R_m\}$ (" $R_1, \dots, R_m$ codes an n -ary relation (operation)").

Then for any T and T' , TQT' implies $T \cong_s T'$.

PROOF. Let S, Σ, s , and Q be as hypothesized. Let C be the category whose objects are countable ω -dimensional locally-finite cylindric algebras with an S -indexed list of constants and whose maps are homomorphisms. Note that if R_1 and R_2 are sets of symbols such that $R_1 \cap R_2 \subseteq S$, then $\text{cyl}_s(Th_s(R_1 \cup R_2)) \cong \text{cyl}_s(Th_s(R_1)) \otimes \text{cyl}_s(Th_s(R_2))$. If T is a theory set of symbols includes S and if σ_1 and σ_2 are sentences of T 's language such that $\sigma_1 = \neg \sigma_2$, then $\text{cyl}_s(T) \cong \text{cyl}_s(T\sigma_1) \oplus \text{cyl}_s(T\sigma_2)$.

Let Q also denote the symmetric relation on algebras of C defined by $\mathcal{U}Q\mathcal{B}$

iff for some theories T and T' $\mathfrak{U} = \text{cyl}_s(T)$, $\mathfrak{B} = \text{cyl}_s(T')$, and TQT' . We show that $\mathbf{Q}$ is Φ -decomposable.

Condition (D1) is satisfied since the minimal subalgebras of related algebras $\text{cyl}_s(T)$ and $\text{cyl}_s(T')$ are $\text{cyl}_s(T|\emptyset)$ and $\text{cyl}_s(T'|\emptyset)$ respectively which, by (C1), are the same.

To verify (D2), suppose $\mathfrak{U} \mathbf{Q} \mathfrak{B}$ and $\mathfrak{E} \rightarrow \mathfrak{U}$ is a monomorphism from a finitely generated algebra. Then there are theories T and T' such that $\mathfrak{U} = \text{cyl}_s(T)$, $\mathfrak{B} = \text{cyl}_s(T')$, and TQT' . By the hypotheses concerning $\mathbf{Q}$, there are sets of symbols R and R' such that $T = Th_s(R)$ and $T' = Th_s(R')$. Since $\mathfrak{E}$ is finitely generated, there is a finite set $F \subseteq R$ such that $\mathfrak{E} \rightarrow \mathfrak{U}$ factors through the injection $\text{cyl}_s(T|F) \rightarrow \text{cyl}_s(T) = \mathfrak{U}$.

Case (i). Suppose $F = \{R\}$ for some symbol $R \in R$ whose type occurs in R' . Let R' be a symbol of R' of the same type as R . Let $\varphi(X) = \text{cyl}_s(T|\{R\}) \otimes X$, let $\mathfrak{U}_1 = \text{cyl}_s(T|R \sim \{R\})$, and let $\mathfrak{B}_1 = \text{cyl}_s(T'|R' \sim \{R'\})$. Then $\varphi(\mathfrak{U}_1) = \text{cyl}_s(T|\{R\}) \otimes \text{cyl}_s(T|R \sim \{R\}) \cong \text{cyl}_s(T) = \mathfrak{U}$; $\varphi(\mathfrak{B}_1) = \text{cyl}_s(T|\{R\}) \otimes \text{cyl}_s(T'|R' \sim \{R'\}) \cong$, since $T|\{R\} = Th_s(\{R\}) \cong_s Th_s(\{R'\}) = T'|\{R'\}$, $\text{cyl}_s(T'|\{R'\}) \otimes \text{cyl}_s(T'|R' \sim \{R'\}) \cong \text{cyl}_s(T') = \mathfrak{B}$; $\mathfrak{U}_1 = \text{cyl}_s(T|R \sim \{R\}) \mathbf{Q} \text{cyl}_s(T'|R' \sim \{R'\}) = \mathfrak{B}_1$ by (C3); and, noting that the minimal subalgebra $\mathfrak{D}_1$ of $\mathfrak{U}_1$ is $\text{cyl}_s(T|\emptyset)$, $\text{cyl}_s(T|\{R\}) \rightarrow \text{cyl}_s(T) = \mathfrak{U}$ factors through $\varphi(\mathfrak{D}_1) \rightarrow \varphi(\mathfrak{U}_1) \cong \mathfrak{U}$, i.e., through $\varphi(\mathfrak{D}_1) = \text{cyl}_s(T|\{R\}) \otimes \text{cyl}_s(T|\emptyset) \rightarrow \text{cyl}_s(T|\{R\}) \otimes \text{cyl}_s(T|R \sim \{R\}) \cong \text{cyl}_s(T) = \mathfrak{U}$.

Case (ii). Suppose $F = \{R\}$ for some n -ary predicate (operation) symbol $R \in R$ whose type does not occur in R' . By (C5), there are symbols $R_1, \dots, R_m$ in R and symbols $R'_1, \dots, R'_m$ in R' of corresponding type and a sentence " $R_1, \dots, R_m$ codes an n -ary relation (operation)" such that $T|\{R\} \cong_s T|\{R_1, \dots, R_m\}$ (" $R_1, \dots, R_m$ codes an n -ary relation (operation)") via some isomorphism π . Let " $R_1, \dots, R_m$ codes a trivial relation (operation)" be the value π assigns to " R is trivial". Let $\sigma_1 = "$ R is nontrivial", $\sigma_2 = "$ R is trivial" $\wedge "$ $R_1, \dots, R_m$ does not code any n -ary relation (operation)", and $\sigma_3 = "$ R is trivial" $\wedge "$ $R_1, \dots, R_m$ codes an n -ary relation (operation)" and let $\sigma'_1 = "$ $R'_1, \dots, R'_m$ codes an n -ary relation (operation)" $\wedge "$ $R'_1, \dots, R'_m$ does not code a trivial relation (operation)", $\sigma'_2 = "$ $R'_1, \dots, R'_m$ does not code an n -ary relation (operation)", and $\sigma'_3 = "$ $R'_1, \dots, R'_m$ codes a trivial relation (operation)". Then σ_1, σ_2 , and σ_3 are pairwise inconsistent and mutually exhaustive as are σ'_1, σ'_2 , and σ'_3 . Also $T|\{R\}(\sigma_1) \cong_s T'|\{R'_1, \dots, R'_m\}(\sigma'_1)$; $T|\{R, R_1, \dots, R_m\}(\sigma_2) \cong_s T'|\{R'_1, \dots, R'_m\}(\sigma'_2)$; $T\sigma_3 \cong_s$, this requires a second glance, $T|\{R \sim \{R_1, \dots, R_m\}\}$; and $T'\sigma'_3 \cong_s T'|\{R' \sim \{R'_1, \dots, R'_m\}\}$. Let

$$\varphi(X_1, X_2, X_3) = (\text{cyl}_s(T|\{R\}\sigma_1) \otimes X_1)$$

$$\oplus (\text{cyl}_s(T|\{R, R_1, \dots, R_m\}\sigma_2) \otimes X_2) \oplus X_3,$$

let $\mathfrak{U}_1 = \text{cyl}_s(T|R \sim \{R\})$, $\mathfrak{U}_2 = \text{cyl}_s(T|R \sim \{R, R_1, \dots, R_m\})$, $\mathfrak{U}_3 = \text{cyl}_s(T|R \sim \{R_1, \dots, R_m\})$, and let

$$\mathfrak{B}_1 = \mathfrak{B}_2 = \mathfrak{B}_3 = \text{cyl}_s(T'|R' \sim \{R'_1, \dots, R'_n\}).$$

Then

$$\begin{aligned} \varphi(\mathfrak{U}_1, \mathfrak{U}_2, \mathfrak{U}_3) &= (\text{cyl}_s(T|\{R\}\sigma_1) \otimes \text{cyl}_s(T|R \sim \{R\})) \\ &\quad \oplus (\text{cyl}_s(T|\{R, R_1, \dots, R_m\}\sigma_2) \\ &\quad \otimes \text{cyl}_s(T|R \sim \{R, R_1, \dots, R_m\})) \\ &\quad \oplus \text{cyl}_s(T|R \sim \{R_1, \dots, R_m\}) \\ &\cong \text{cyl}_s(T\sigma_1) \oplus \text{cyl}_s(T\sigma_2) \oplus \text{cyl}_s(T\sigma_3) \cong \text{cyl}_s(T) = \mathfrak{U}. \end{aligned}$$

Similarly, $\varphi(\mathfrak{B}_1, \mathfrak{B}_2, \mathfrak{B}_3) \cong \mathfrak{B}$. By (C3), $\mathfrak{U}_1 \mathbf{Q} \mathfrak{B}_1$, $\mathfrak{U}_2 \mathbf{Q} \mathfrak{B}_2$, and $\mathfrak{U}_3 \mathbf{Q} \mathfrak{B}_3$. Finally, the minimal subalgebras of $\mathfrak{U}_1$, $\mathfrak{U}_2$, and $\mathfrak{U}_3$ are $\mathfrak{D}_1 = \mathfrak{D}_2 = \mathfrak{D}_3 = \text{cyl}_s(T|\emptyset)$, and $\text{cyl}_s(T|\{R\}) \rightarrow \text{cyl}_s(T) = \mathfrak{U}$ factors through

$$\varphi(\mathfrak{D}_1, \mathfrak{D}_2, \mathfrak{D}_3) \rightarrow \varphi(\mathfrak{U}_1, \mathfrak{U}_2, \mathfrak{U}_3) = \mathfrak{U}$$

which is

$$\begin{aligned} &(\text{cyl}_s(T|\{R\}\sigma_1) \otimes \text{cyl}_s(T|\emptyset)) \\ &\quad \oplus (\text{cyl}_s(T|\{R, R_1, \dots, R_m\}\sigma_2) \otimes \text{cyl}_s(T|\emptyset)) \\ &\quad \oplus \text{cyl}_s(T|\emptyset) \rightarrow \text{cyl}_s(T|\sigma_1) \oplus \text{cyl}_s(T|\sigma_2) \oplus \text{cyl}_s(T|\sigma_3) \cong \mathfrak{U} \end{aligned}$$

since if σ_1 holds R is in $\{R\}$, if σ_2 holds R is in $\{R, R_1, \dots, R_m\}$, and if σ_3 holds R is trivial and hence definable from S .

Case (iii). Suppose $\bar{F} = \{R_1, \dots, R_n\}$. By (C2) there are $R_1, \dots, R_n \subseteq R$ and $R'_1, \dots, R'_n \subseteq R'$ such that $R_i \in R_i$, $R_i \cap R_j = R'_i \cap R'_j = S$ for $i \neq j$, $R = R_1 \cup \dots \cup R_n$, $R' = R'_1 \cup \dots \cup R'_n$, and $T|R_i \mathbf{Q} T'|R'_i$. Let $\mathfrak{U}_i = \text{cyl}_s(T|R_i)$ and $\mathfrak{B}_i = \text{cyl}_s(T'|R'_i)$. Then $\mathfrak{U}_i \mathbf{Q} \mathfrak{B}_i$. By Cases (i) and (ii), there are $\varphi_i \in \Phi$, sequences $\bar{\mathfrak{U}}_i = \langle \mathfrak{U}_i^1, \mathfrak{U}_i^2, \dots \rangle$ and $\bar{\mathfrak{B}}_i = \langle \mathfrak{B}_i^1, \mathfrak{B}_i^2, \dots \rangle$, and isomorphisms $\varphi_i(\bar{\mathfrak{U}}_i) \cong \mathfrak{U}_i$ and $\varphi_i(\bar{\mathfrak{B}}_i) \cong \mathfrak{B}_i$ such that $\mathfrak{U}_i^j \mathbf{Q} \mathfrak{B}_i^j$ and $\text{cyl}_s(T|\{R_i\}) \rightarrow \text{cyl}_s(T|R_i) = \mathfrak{U}_i$ factors through $\varphi_i(\bar{\mathfrak{D}}_i) \rightarrow \varphi_i(\bar{\mathfrak{U}}_i) \cong \mathfrak{U}_i$ where $\bar{\mathfrak{D}}_i$ is the minimal subalgebra sequence of $\bar{\mathfrak{U}}_i$. Let φ be $\varphi_1 \otimes \varphi_2 \otimes \dots \otimes \varphi_n$, and let $\bar{\mathfrak{U}} = \bar{\mathfrak{U}}_1 \frown \bar{\mathfrak{U}}_2 \frown \dots \frown \bar{\mathfrak{U}}_n$ and $\bar{\mathfrak{B}} = \bar{\mathfrak{B}}_1 \frown \bar{\mathfrak{B}}_2 \frown \dots \frown \bar{\mathfrak{B}}_n$ where $\frown$ is concatenation. Then $\varphi(\bar{\mathfrak{U}}) \cong \varphi_1(\bar{\mathfrak{U}}_1) \otimes \dots \otimes \varphi_n(\bar{\mathfrak{U}}_n) \cong \mathfrak{U}_1 \otimes \dots \otimes \mathfrak{U}_n \cong \mathfrak{U}$ and likewise $\varphi(\bar{\mathfrak{B}}) \cong \mathfrak{B}$. Also $\mathfrak{U}_i^j \mathbf{Q} \mathfrak{B}_i^j$ for each $\mathfrak{U}_i^j$ and $\mathfrak{B}_i^j$ of $\bar{\mathfrak{U}}$ and $\bar{\mathfrak{B}}$. Finally, $\text{cyl}_s(T|\{R_1, \dots, R_n\}) \rightarrow \text{cyl}_s(T) = \mathfrak{U}$ factors through

$$\begin{aligned} &\text{cyl}_s(T|\{R_1\}) \otimes \dots \otimes \text{cyl}_s(T|\{R_n\}) \\ &\rightarrow \text{cyl}_s(T|R_1) \otimes \dots \otimes \text{cyl}_s(T|R_n) \cong \text{cyl}_s(T) = \mathfrak{U} \end{aligned}$$

which factors through $\varphi(\bar{\mathfrak{D}}) \rightarrow \varphi(\bar{\mathfrak{U}}) \cong \mathfrak{U}$, i.e., through

$$\begin{aligned}\varphi(\bar{\mathfrak{D}}) &= \varphi(\bar{\mathfrak{D}}_1) \otimes \cdots \otimes \varphi_n(\bar{\mathfrak{D}}_n) \rightarrow \varphi_1(\bar{\mathfrak{U}}_1) \otimes \cdots \otimes \varphi_n(\bar{\mathfrak{U}}_n) \\ &\cong \mathfrak{U}_1 \otimes \cdots \otimes \mathfrak{U}_n \cong \mathfrak{U}.\end{aligned}$$

Hence $\mathbf{Q}$ is Φ -decomposable and by the preceding lemma TQT' implies $T \cong_s T'$.

Basic isomorphisms.

LEMMA 9. *For any sets $\mathcal{R}$ and $\mathcal{R}'$ of symbols with no constants, if there are arbitrarily large integers such that both $\mathcal{R}$ and $\mathcal{R}'$ have predicate symbols of that type or if $\mathcal{R}$ and $\mathcal{R}'$ have the same unboundedness, the unboundedness is an integer, and no symbols of $\mathcal{R}$ or $\mathcal{R}'$ have type larger than the unboundedness; then $Th(\mathcal{R}) \cong Th(\mathcal{R}')$.*

PROOF. Let $\mathbf{Q}$ be the relation such that TQT' iff $T = Th(\mathcal{R})$ and $T' = Th(\mathcal{R}')$ for some $\mathcal{R}$ and $\mathcal{R}'$ satisfying the hypotheses of the lemma. Suppose $T = Th(\mathcal{R})$ and $T' = Th(\mathcal{R}')$ and TQT' . Then conditions (C1)–(C4) of the Construction theorem hold. To show (C5), suppose $R \in \mathcal{R}$ is a symbol of a type which does not occur in $\mathcal{R}'$. Then the hypotheses imply there is a predicate symbol $R_0 \in \mathcal{R}$ of a type m which is larger than that of R and which does occur in $\mathcal{R}'$.

Case (i). R is an n -ary predicate symbol. Let “ R_0 codes an n -ary relation” be the sentence

$$\begin{aligned}(\forall x_1 \cdots x_n) (\forall x_{n+1} \cdots x_m x'_{n+1} \cdots x'_m) \\ (Rx_1 \cdots x_n x_{n+1} \cdots x_m \iff Rx_1 \cdots x_n x'_{n+1} \cdots x'_m).\end{aligned}$$

Then $T|\{R\} \cong T|\{R_0\}$ (“ R_0 codes an n -ary relation”).

Case (ii). R is an n -ary operation symbol. Let “ R_0 codes an n -ary operation” be the sentence

$$(\forall x_1 \cdots x_n) (\exists y) (\forall x_{n+1} \cdots x_m) (Rx_1 \cdots x_n x_{n+1} \cdots x_m \iff x_{n+1} = y).$$

Then $T|\{R\} \cong T|\{R_0\}$ (“ R_0 codes an n -ary operation”).

Let $\exists \geq 2$ be the sentence $(\exists xy) (x \neq y)$ and let “ R is irreflexive” be $(\forall x) (\neg Rx \cdots x)$.

LEMMA 10. *For any sets $\mathcal{R}$ and $\mathcal{R}'$ of symbols with no constants, if there are arbitrarily large half integers such that $\mathcal{R}$ and $\mathcal{R}'$ have operation symbols of that type or if $\mathcal{R}$ and $\mathcal{R}'$ have the same unboundedness, the unboundedness is a half integer $\geq 2\frac{1}{2}$, and no symbols of $\mathcal{R}$ and $\mathcal{R}'$ have type larger than the unboundedness, then $Th(\mathcal{R}) (\exists \geq 2) \cong Th(\mathcal{R}') (\exists \geq 2)$.*

PROOF. Let $s = \langle \emptyset, \{\exists \geq 2\} \rangle$. Let $\mathbf{Q}$ be the relation such that TQT' iff $T = Th_s(\mathcal{R})$ and $T' = Th_s(\mathcal{R}')$ for some $\mathcal{R}$ and $\mathcal{R}'$ satisfying the hypotheses of

the lemma. Suppose $T = Th_s(R)$ and $T' = Th_s(R')$ and TQT' . Then conditions (C1)–(C4) of the Construction theorem hold. To show (C5) suppose $R \in \mathcal{R}$ is a symbol of a type which does not occur in $\mathcal{R}'$. Then the hypotheses imply there are operation symbols R_0 and R_1 in $\mathcal{R}$ of types $m + \frac{1}{2}$, $m \geq 2$, and $k + \frac{1}{2}$, $k \geq 2$, which are larger than that of R and which do occur in $\mathcal{R}'$.

Case (i). R is an n -ary operation symbol. Let “ R_0 codes an n -ary operation” be the sentence

$$(\forall x_1 \cdots x_n) (\forall x_{n+1} \cdots x_m x'_{n+1} \cdots x'_m) \\ (Rx_1 \cdots x_n x_{n+1} \cdots x_m = Rx_1 \cdots x_n x'_{n+1} \cdots x'_m).$$

Then $T|\{R\} \cong T|\{R_0\}$ (“ R_0 codes an n -ary operation”).

Case (ii). R is an n -ary predicate symbol, $n \geq 1$. Let “ R_0 codes an irreflexive n -ary relation” be “ R_0 codes an n -ary operation” $\wedge$

$$(\forall x_1 \cdots x_n x_{n+1} \cdots x_m) \\ \left(Rx_1 \cdots x_m = x_1 \vee \bigvee_{1 \leq i \leq n} (x_1 = x_2 = \cdots = x_i \right. \\ \left. \neq x_{i+1} \wedge Rx_1 \cdots x_m = x_{i+1}) \right),$$

and let “ R_1 codes a unary relation” be

$$(\forall x_1) [(\forall x_2 \cdots x_k (Rx_1 \cdots x_k = x_1)) \vee (\forall x_2 \cdots x_k (Rx_1 \cdots x_k = x_2))].$$

Then for any unary predicate symbol $U \notin \mathcal{R}$, $T|\{R\} \cong T|\{R\}$ (“ R is irreflexive”) “ $\otimes$ ” $(T|\{R \sim \{R\}\} \cup \{U\}) \cong T|\{R_0, R_1\}$ (“ R_0 codes an irreflexive n -ary relation” $\wedge$ “ R_1 codes a unary relation”).

Case (iii). R is a 0-ary predicate. Let “ R_0 codes a 0-ary relation” be $(\forall x_1 \cdots x_m (Rx_1 \cdots x_m = x_1)) \vee (\forall x_1 \cdots x_m (Rx_1 \cdots x_m = x_2))$. Then $T|\{R\} \cong T|\{R_0\}$ (“ R_0 codes a 0-ary predicate”).

LEMMA 11. For any sets of symbols $\mathcal{R}$ and $\mathcal{R}'$ with at least one constant and infinitely many symbols of type ≥ 1 , if $\mathcal{R}$ and $\mathcal{R}'$ have exactly the same number of symbols of each type other than $\frac{1}{2}$, then $Th(\mathcal{R}) \cong Th(\mathcal{R}')$.

PROOF. We may assume without loss of generality that $\mathcal{R}$ and $\mathcal{R}'$ contain a common constant a . Let $s = \langle \{a\}, \emptyset \rangle$. Let Q be the relation such that TQT' iff $T = Th_s(R)$ and $T' = Th_s(R')$ for some $\mathcal{R}$ and $\mathcal{R}'$ satisfying the hypotheses of the lemma and the assumption above. Suppose $T = Th_s(R)$ and $T' = Th_s(R')$ and TQT' . Then conditions (C1)–(C4) of the Construction theorem hold. To show (C5) suppose $R \in \mathcal{R}$ is a symbol of a type which does not occur in $\mathcal{R}'$. Then the hypotheses imply R is some constant b and there is a predicate or operation symbol R_0 of $\mathcal{R}$ of a type m or $m + \frac{1}{2}$, $m \geq 1$, which also occurs in $\mathcal{R}'$.

If R_0 is a predicate symbol, let “ R_0 codes a constant” be $(\exists y) (\forall x_1 \cdots x_m) (Rx_1 \cdots x_m \iff x_1 = y)$. If R_0 is an operation symbol, let “ R_0 codes a

constant" be $(\exists y) (\forall x_1 \cdots x_m) (Rx_1 \cdots x_m = y)$. Then $T| \{b\} \cong_s T| \{R_0\}$ (" R_0 codes a constant").

LEMMA 12. *For any set of symbols R and R' with at least one constant and infinitely many symbols of type ≥ 1 , if R and R' have exactly the same number of symbols of each type other than 0 then $Th(R) (\exists \geq 2) \cong Th(R') (\exists \geq 2)$.*

PROOF. We may assume that R and R' contain a common constant a . Let $s = \langle \{a\}, \{\exists \geq 2\} \rangle$. Let Q be the relation such that TQT' iff $T = Th_s(R)$ and $T' = Th_s(R')$ for some R and R' satisfying the hypotheses of the lemma and the assumption above. Suppose $T = Th_s(R)$ and $T' = Th_s(R')$ and TQT' . Then conditions (C1)–(C4) of the Construction theorem hold. To show (C5) suppose $R \in R$ is a symbol of a type which does not occur in R' . Then the hypotheses imply R is some 0-ary predicate P and there is a predicate or operation symbol R_0 of R of a type m or $m + \frac{1}{2}$, $m \geq 1$, which also occurs in R' .

If R_0 is a predicate symbol, let " R_0 codes a 0-ary predicate" be $(\forall x_1 \cdots x_m (R_0 x_1 \cdots x_m)) \vee (\forall x_1 \cdots x_m (\neg R_0 x_1 \cdots x_m))$. If R_0 is an operation symbol, let " R_0 codes a 0-ary predicate" be $(\forall x_1 \cdots x_m (R_0 x_1 \cdots x_m = x_1)) \vee (\forall x_1 \cdots x_m (R_0 x_1 \cdots x_m = a))$. Then $T| \{P\} \cong_s T| \{R_0\}$ (" R_0 codes a 0-ary predicate").

LEMMA 13. *For any sets of symbols R and R' with at least one constant, infinitely many symbols of type ≥ 1 , and infinitely many of type 0, if R and R' have exactly the same number of symbols of each type other than 1, then $Th(R) (\exists \geq 2) \cong Th(R') (\exists \geq 2)$.*

PROOF. We may assume without loss of generality that R and R' contain a common constant a . Let $s = \langle \{a\}, \{\exists \geq 2\} \rangle$. Let Q be the relation such that TQT' iff $T = Th_s(R)$ and $T' = Th_s(R')$ for some R and R' satisfying the hypotheses of the lemma and the assumption above. Suppose $T = Th_s(R)$ and $T' = Th_s(R')$ and TQT' . Then conditions (C1)–(C4) of the Construction theorem hold. To show (C5) suppose $R \in R$ is a symbol of a type which does not occur in R' . Then the hypotheses imply R is some unary predicate U and there is 0-ary predicate symbol $P \in R$ and a predicate or operation symbol $R_0 \in R$ of a type m or $m + \frac{1}{2}$, $m \geq 1$ which also occurs in R' .

If R_0 is a predicate symbol, let " R_0, P codes a unary predicate" be

$$(\forall x_1) (\forall x_2 \cdots x_m x'_2 \cdots x'_m) (Rx_1 x_2 \cdots x_m \iff Rx_1 x'_2 \cdots x'_m).$$

If R_0 is an operation symbol, let " R_0, P codes a unary predicate" be $(\forall x_1 \cdots x_m) (Rx_1 \cdots x_m = x_1 \vee Rx_1 \cdots x_m = a)$. The idea is that $x \in U$ is determined by R_0 for $x \neq a$ and by P for $x = a$. Then $T| \{U\} \cong_s T| \{R_0, P\}$ (" R_0, P codes a unary predicate").

Let $\exists!1$ be the sentence $(\forall xy)(x = y)$.

LEMMA 14. *For any sets R and R' of symbols, if R and R' have the same number of predicate symbols, then $Th(R) (\exists!1) \cong Th(R') (\exists!1)$.*

PROOF. On a set with exactly one element, an operation symbol has exactly one possible interpretation and a predicate symbol exactly two. Thus all operation symbols are definable and eliminable and all predicate symbols can be replaced by 0-ary predicates. Hence the lemma.

Classification theorem.

THEOREM 15 (CLASSIFICATION THEOREM). *The following properties and functions are a complete set of invariants for countable languages under cylindric isomorphisms.*

- (1) *The reduced type sequence.*
- (2) *The existence of a constant.*
- (3) *The number of predicate symbols.*
- (4) *The number of unary predicate symbols in case the language's only other symbols are for 0-ary predicates and unary operations.*

PROOF. The four properties have been shown to be invariants. Suppose T and T' are the theories of two languages L and L' with symbols R and R' which are the same with respect to the invariants. We must show that T and T' are cylindrically isomorphic.

Since L and L' have the same reduced type sequence, they have the same unboundedness. If the unboundedness is undefined, R and R' are finite and so the reduced type sequence is the full type sequence and hence $T \cong T'$. Suppose the unboundedness is defined. Let R_0 and R'_0 be the set of symbols of R and R' respectively whose type is less than or equal to the unboundedness. Then $R \sim R_0$ and $R' \sim R'_0$ are finite and have the same type sequence. Hence $T|(R \sim R_0) \cong T'|(R' \sim R'_0)$. Thus to show $T \cong T'$ it suffices to show $T|R_0 \cong T'|R'_0$. Consequently, we may henceforth assume R and R' have no symbols of type greater than the unboundedness.

If the unboundedness is 0, the reduced type sequence is the full sequence and hence $T \cong T'$. If the unboundedness is $\frac{1}{2}$, the reduced type sequence and the number of predicate symbols determine the full type sequence and hence $T \cong T'$. Henceforth, assume the unboundedness of R and R' is ≥ 1 .

Suppose R and R' have constants a and a' respectively. Let R_0 and R'_0 be R and R' respectively minus all constants except a and a' . By Lemma 11, $T \cong T|R_0$ and $T' \cong T'|R'_0$. Hence to show $T \cong T'$, it suffices to show $T|R_0 \cong T'|R'_0$. Henceforth assume R and R' have no constants or exactly one constant each, a and a' respectively.

Since R and R' have the same number of predicate symbols, $T(\exists!1) \cong T'(\exists!1)$ by Lemma 14. It remains to be shown that $T(\exists \geq 2) \cong T'(\exists \geq 2)$.

Suppose the unboundedness is $\neq 1\frac{1}{2}$ and $\neq \omega$. If there are no constants, then $T(\exists \geq 2) \cong T'(\exists \geq 2)$ by Lemma 9 or 10. If there are unique constants a and a' , then $T(\exists \geq 2) \upharpoonright \{a\} \cong T'(\exists \geq 2) \upharpoonright \{a'\}$ and, by Lemma 9 or 10, $T(\exists \geq 2) \upharpoonright (R \sim \{a\}) \cong T(\exists \geq 2) \upharpoonright (R' \sim \{a'\})$. Hence $T(\exists \geq 2) \cong T'(\exists \geq 2)$.

Suppose the unboundedness is ω . Then, in the case of no constants, $T(\exists \geq 2) = Th(R) (\exists \geq 2) \cong$, by Lemma 9 or 10, $Th(R \cup R') (\exists \geq 2) \cong Th(R') (\exists \geq 2) = T'(\exists \geq 2)$. The case with a unique constant is handled similarly as in the previous paragraph.

Suppose the unboundedness is $1\frac{1}{2}$ and there are no constants. If the number of unary predicates in R and/or R' is finite, then the reduced type sequence, the number of predicate symbols, and the number of unary predicate symbols determine the full type sequence and hence $T(\exists \geq 2) \cong T'(\exists \geq 2)$. If there are infinitely many unary predicates, let R_0 and R'_0 be the set of unary operation symbols of R and R' respectively. Then $T(\exists \geq 2) \upharpoonright R_0 \cong T'(\exists \geq 2) \upharpoonright R'_0$ is clear and $T(\exists \geq 2) \upharpoonright (R \sim R_0) \cong T'(\exists \geq 2) \upharpoonright (R' \sim R'_0)$ follows from Lemma 9. Hence $T(\exists \geq 2) \cong T'(\exists \geq 2)$.

Suppose the unboundedness is $1\frac{1}{2}$ and there are unique constants a and a' . Let U be a countably infinite set of propositional letters disjoint from R and R' . Then $T(\exists \geq 2) = Th(R) (\exists \geq 2) \cong$, by Lemma 12, $Th(R \cup U) (\exists \geq 2) \cong$, by Lemma 13, $Th(R' \cup U) (\exists \geq 2) \cong Th(R') (\exists \geq 2) = T(\exists \geq 2)$.

Hence $T \cong T'$.

While we have considered only countable languages, the generalization of the Classification theorem to languages of arbitrary cardinality is not too difficult.

The corresponding classification problem for the Boolean algebras of sentences of countable languages has been solved by Hanf and his student, Simons. In this case all languages of infinite similarity type have isomorphic Boolean algebras. So do all languages of finite similarity type with at least one symbol of type ≥ 2 or at least two symbols of type $1\frac{1}{2}$. So do all languages of finite similarity type with no symbols of type ≥ 2 and exactly one of type $1\frac{1}{2}$. The remaining cases are all nonisomorphic.

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